

Construction of Uniformly Distributed Linear Recurring Sequences Over Dedekind Domains

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Pseudo Random Number Generation

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e.g. real sequences:

$$\alpha \cdot n \bmod 1$$

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Sequences can be

- theoretical
e.g. real sequences:

$$\alpha \cdot n \bmod 1$$

- practical e.g. sequences over finite structures:

$$n \bmod M.$$

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- easy computation

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- good statistical properties, e.g.

 - good approximation to the distribution

 - low correlation of consecutive elements

Main result

If D is Dedekind domain, $\mathcal{I} \subseteq D$ is a particular prime ideal, u is a linear recurring sequence in D

- with a recurrence relation satisfying some simple condition,
- and having maximal period length $\pmod{\mathcal{I}}$,

then u is uniformly distributed $\pmod{\mathcal{I}^s}$ for all positive integers s .

Dedekind domains

D : a Dedekind domain, $\mathcal{I} \subseteq D$: an ideal

Definition (Norm)

$$\text{Norm: } \mathbf{N}(\mathcal{I}) = |D/\mathcal{I}|$$

$$\text{Finite norm: } \mathbf{N}(\mathcal{I}) < \infty$$

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Finite norm: $\mathbf{N}(\mathcal{I}) < \infty$

$u \in D^\infty$ a sequence in D

Definition (Uniform distribution)

\mathcal{I} has finite norm: u is **uniformly distributed** modulo \mathcal{I} , if

$$\lim_{m \rightarrow \infty} \frac{1}{m} |\{n < m \mid u_n \equiv r \pmod{\mathcal{I}}\}| = \frac{1}{\mathbf{N}(\mathcal{I})} \quad \text{for all } r \in D$$

Lemma (Semi GCD domain)

$Q_1, Q_2 \in D[x]$, Q_1 is monic.

Then there exist unique $\gcd(Q_1, Q_2)$ and $\text{lcm}(Q_1, Q_2)$.

The gcd is monic.

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If both Q_1, Q_2 are monic: lcm is monic.

Definition (Generating function)

$u \in D^\infty$, $G_u \in D[[x]]$.

Generating function of u :
$$G_u(x) = \sum_{n=0}^{\infty} u_n \cdot x^n$$

Linear Recurring Sequences

D : a Dedekind domain, $u \in D^\infty$ a sequence in D

Definition (LRS)

$a_0, \dots, a_{d-1} \in D$, u satisfies the **recurrence relation**

$$u_{n+d} = a_{d-1}u_{n+d-1} + \dots + a_0u_n \quad n = 0, 1, \dots$$

u : **linear recurring sequence**

a_0, \dots, a_{d-1} : **coefficients**

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$$P \in D[x]$$

Definition (Characteristic polynomial)

characteristic polynomial: $P(x) = x^d - a_{d-1}x^{d-1} - \dots - a_0$.

Lemma (LRS and Generating function)

$u \in D^\infty$, $G_u \in D[[x]]$ the generating function of u .

u is a LRS $\iff \exists P \in D[x]$, s.t. $P^* \cdot G_u \in D[x]$.
 P^* is the reciprocal polynomial of P .

Remark

Technically: P is a characteristic polynomial of u .

Lemma

$P \in D[x]$ is monic, $d = \deg(P)$,
 $\mathcal{U}(D, P) = \{u \mid u \in D^\infty, \deg(P^* \cdot G_u) < d\}$

Then $\mathcal{U}(D, P) \cong D[x]/x^d D[x] \cong D^d$.

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Lemma

$P, Q \in D[x]$ are monic, then

$$\mathcal{U}(D, P) \subseteq \mathcal{U}(D, P \cdot Q)$$

$$\mathcal{U}(D, P) \cap \mathcal{U}(D, Q) = \mathcal{U}(D, \gcd(P, Q))$$

Lemma (Linear combination of LRSs)

$$\mathcal{U}(D) = \bigcup_{\substack{P \in D[x] \\ \text{monic}}} \mathcal{U}(D, P) \text{ is a module.}$$

Module structure of LRSs

Lemma (Linear combination of LRSs)

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Definition (Linear complexity)

$u \in D^\infty$, s.t. $G_u \in \mathcal{U}(D)$.

The **linear complexity** of u :

$$d(u) = \min \{ \dim(\mathcal{U}(D, P)) \mid P \in D[x] \text{ monic}, u \in \mathcal{U}(D, P) \}$$

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Lemma (Minimal characteristic polynomial)

If u is a LRS, then there exists a unique $P \in D[x]$ monic polynomial, s.t.

$$d(u) = \dim(\mathcal{U}(D, P))$$

Definition

$u \in D^\infty$ is **(ultimately) periodic**, if $\exists \varrho \in \mathbb{Z}^+$, s.t.

$$G_u \in \mathcal{U}(D, x^\varrho - 1)$$

ϱ is a **period length**;

The smallest such a ϱ is the **minimal period length**, denoted by $\varrho(u)$.

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Lemma

$u \in D^\infty$ is a LRS, $\mathcal{I} \subseteq D$ is an ideal with finite norm.

Then u is periodic mod \mathcal{I} .

Definition (Impulse response sequence (IRS))

$u \in D^\infty$, s.t. $G_u \in \mathcal{U}(D)$, $d = d(u)$.

u is an **impulse response sequence**, if

$$u_0 = \cdots = u_{d-2} = 0, \quad u_{d-1} = 1 .$$

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If u is an IRS, then

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Lemma (Periodicity of LRS)

$u \in D^\infty$ is a LRS, $\mathcal{I} \subseteq D$ is an ideal with finite norm.

Then u is periodic mod \mathcal{I} .

Lemma (Maximal period length)

\mathbb{F} is a finite field, $P \in \mathbb{F}[x]$ monic, and $u, v \in \mathbb{F}^\infty$, s.t $G_u, G_v \in \mathcal{U}(\mathbb{F}, P)$.

Then u, v are periodic, and if u is an IRS, then $\varrho(v) | \varrho(u)$.

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Lemma (Shifted sequence)

Let $Q, P \in D[x]$, s.t. Q is monic and $P(x) = (x - 1)^2 Q(x)$,
 $a \in D$ and $u, v \in D^\infty$, s.t. $v_n = u_n + a$ for all $n \geq 0$.

Then $G_u \in \mathcal{U}(D, P) \implies G_v \in \mathcal{U}(D, P)$

Theorem (Uniform distribution)

$P, P', Q \in D[x]$ monic, s.t

$$P(x) = (x - 1)^2 Q(x) \text{ and } P'(x) = (x - 1)Q(x)$$

$u, v \in D^\infty$, s.t. $G_u, G_v \in \mathcal{U}(D, P)$, and v is an IRS,

$\mathcal{I} \subset D$ is a prime ideal with $\mathbf{N}(\mathcal{I}) = p$,

Q is irreducible mod \mathcal{I} ,

$\varrho(\mathcal{I}, u)$ and $\varrho(\mathcal{I}, v)$ are the period lengths mod \mathcal{I} .

If $\varrho(\mathcal{I}, u) = \varrho(\mathcal{I}, v)$, then u uniformly distributed mod \mathcal{I} .

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If $\varrho(\mathcal{I}, u) = \varrho(\mathcal{I}, v)$, then u uniformly distributed mod \mathcal{I} .

The proof is based on the observation of the structure of $\mathcal{U}(D/\mathcal{I}, P)/\mathcal{U}(D/\mathcal{I}, P')$.

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$p \in \mathbb{N}$ is a prime,

$\mathcal{I} \subset D$ is a prime ideal with $\mathbf{N}(\mathcal{I}) = p$,

Q is irreducible mod \mathcal{I} ,

$s \in \mathbb{Z}^+$, and $\varrho(\mathcal{I}^s, u)$ and $\varrho(\mathcal{I}^s, v)$ are the period lengths mod \mathcal{I}^s .

If $\varrho(\mathcal{I}, u) = \varrho(\mathcal{I}, v)$, then $\varrho(\mathcal{I}^{s+1}, u) = p \cdot \varrho(\mathcal{I}^s, u)$.

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 $p \in \mathbb{N}$ is a prime,
 $\mathcal{I} \subset D$ is a prime ideal with $\mathbf{N}(\mathcal{I}) = p$,
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If $\varrho(\mathcal{I}, u) = \varrho(\mathcal{I}, v)$, then

u is uniformly distributed mod \mathcal{I}^s , for all $s \in \mathbb{Z}^+$.

Furthermore,

$$\varrho(\mathcal{I}^s, u) = \text{ord}(Q) \cdot p^s.$$

$\text{ord}(Q) | p^{\deg(Q)} - 1$, is the order of Q .

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